**Theorem 1.** Suppose we have the following LP in standard form

(P) 
$$\min \mathbf{c}^T \mathbf{x} \ s.t \ \mathbf{A} \mathbf{x} = \mathbf{b} \ and \ \mathbf{x} \ge \mathbf{0},$$

and that  $B = (j_1, \ldots, j_m)$  is a basis.

- (a) The vector  $\mathbf{x}$  defined by  $\mathbf{x}_B = \mathbf{A}_B^{-1}\mathbf{b}$  and  $x_j = 0$  for every nonbasic variable  $x_j$  is the unique basic solution associated with B. In particular, B is a feasible basis when  $\mathbf{A}_B^{-1}\mathbf{b} \geq \mathbf{0}$ .
- (b) If B is a feasible basis and the relative cost vector is non-negative (i.e.  $\mathbf{\bar{c}}^T \geq \mathbf{0}^T$ ), then x is an optimum for (P).
- (c) If B is a feasible basis and there exists a variable  $x_{j_0}$  such that the relative cost of  $x_{j_0}$  is negative (i.e.  $\overline{c}_{j_0} < 0$ ) and  $\mathbf{A}_B^{-1}\mathbf{A}_{j_0} \leq \mathbf{0}$ , then (P) is unbounded.

*Proof.* We begin by proving (a). By the definition,  $\mathbf{x}$  is a basic solution associated with the basis B if and only if  $\mathbf{A}\mathbf{x} = \mathbf{b}$  and  $x_j = 0$  for every nonbasic variable  $x_j$ . Using these two facts we have that

$$\mathbf{b} = \mathbf{A}\mathbf{x}$$
  
=  $\mathbf{A}_1 x_1 + \dots + \mathbf{A}_n x_n$   
=  $\mathbf{A}_{j_1} x_{j_1} + \dots + \mathbf{A}_{j_m} x_{j_m}$   
=  $\left( \mathbf{A}_{j_1} \mid \dots \mid \mathbf{A}_{j_m} \right) \begin{pmatrix} x_{j_1} \\ \vdots \\ x_{j_m} \end{pmatrix}$   
=  $\mathbf{A}_B \mathbf{x}_B$ .

Furthermore, because B is a basis,  $\mathbf{A}_B$  is an invertible matrix. This implies that  $\mathbf{A}_B \mathbf{x}_B = \mathbf{b}$  if and only if  $\mathbf{x}_B = \mathbf{A}_B^{-1}\mathbf{b}$ . Therefore,  $\mathbf{x}$  is a basic solution associated with the basis B if and only if  $\mathbf{x}_B = \mathbf{A}_B^{-1}\mathbf{b}$  and  $x_j = 0$  for every nonbasic variable  $x_j$ . This proves (a).

To prove (b), first note that, since  $x_j = 0$  for every nonbasic variable  $x_j$ ,

(\*)  

$$\mathbf{c}^{T}\mathbf{x} = c_{1}x_{1} + \dots c_{n}x_{n}$$

$$= c_{j_{1}}x_{j_{1}} + \dots c_{j_{m}}x_{j_{m}}$$

$$= \left( c_{j_{1}}, \dots, c_{j_{m}} \right) \begin{pmatrix} x_{j_{1}} \\ \vdots \\ x_{j_{m}} \end{pmatrix}$$

$$= \mathbf{c}_{B}^{T}\mathbf{x}_{B}.$$

Now let  $\mathbf{x}'$  be an arbitrary feasible solution to (P). We will show that  $\mathbf{c}^T \mathbf{x}' \ge \mathbf{c}^T \mathbf{x}$  which will prove (b). By the definition of a feasible solution, we have that  $\mathbf{A}\mathbf{x}' = \mathbf{b}$  and  $\mathbf{x}' \ge \mathbf{0}$ . Using these two facts, the assumption that  $\mathbf{\bar{c}}^T \ge \mathbf{0}^T$  and that fact that, by (a),  $\mathbf{x}_B = \mathbf{A}_B^{-1}b$ , we can

compute that

$$0 \leq \overline{\mathbf{c}}^T \mathbf{x}'$$
  
=  $(\mathbf{c}^T - \mathbf{c}_B^T \mathbf{A}_B^{-1} \mathbf{A}) \mathbf{x}'$   
=  $\mathbf{c}^T \mathbf{x}' - (\mathbf{c}_B^T \mathbf{A}_B^{-1} \mathbf{A}) \mathbf{x}'$   
=  $\mathbf{c}^T \mathbf{x}' - \mathbf{c}_B^T \mathbf{A}_B^{-1} (\mathbf{A} \mathbf{x}')$   
=  $\mathbf{c}^T \mathbf{x}' - \mathbf{c}_B^T \mathbf{A}_B^{-1} \mathbf{b}$   
=  $\mathbf{c}^T \mathbf{x}' - \mathbf{c}_B^T \mathbf{x}_B$ ,

This, with (\*), gives us that

$$\mathbf{c}^T \mathbf{x}' \geq \mathbf{c}_B^T \mathbf{x}_B = \mathbf{c}^T \mathbf{x},$$

which proves (b).

Finally, we prove (c). First, recall that  $\mathbf{x}$  is a feasible solution. We will construct  $\mathbf{y}$  such that Ay = 0,  $y \ge 0$  and  $c^T y < 0$ . This with question 4 on homework assignment 1, will prove (c). First note that  $x_{i_0}$  cannot be a basic variable, because for every  $i \in [m]$ , the relative cost of the basic variable  $x_{j_i}$  is

$$\bar{c}_{j_i} = c_{j_i} - \mathbf{c}_B^T \mathbf{A}_B^{-1} \mathbf{A}_{j_i} = c_{j_i} - \mathbf{c}_B^T (\mathbf{A}_B^{-1} \mathbf{A}_{j_i}) = c_{j_i} - (c_{j_1}, \dots, c_{j_m}) \mathbf{e}_i = c_{j_i} - c_{j_i} = 0,$$

and  $\bar{c}_{j_0} < 0$  by assumption. Here we used that fact that  $\mathbf{A}_{j_i}$  is the *i*th column of  $\mathbf{A}_B$ , so the product  $\mathbf{A}_B^{-1} \mathbf{A}_{j_i}$  must be the *i*th standard basis vector which we denote by  $\mathbf{e}_i$ . Therefore, we can set the variables  $y_{j_1}, \ldots, y_{j_m}$  of **y** by  $\mathbf{y}_B = -\mathbf{A}_B^{-1}\mathbf{A}_{j_0}$  and set  $y_{j_0} = 1$  without conflict. We set the entries in **y** that we have not specified to 0, i.e. we set  $y_j = 0$  for every nonbasic variable  $x_j$  that is not  $x_{j_0}$ . Note that  $\mathbf{y} \geq \mathbf{0}$  because  $\mathbf{A}_B^{-1}\mathbf{A}_{j_0} \leq \mathbf{0}$  by assumption. We now show that  $\mathbf{y}$  is in the null space of  $\mathbf{A}$  with the following computation

$$Ay = A_B y_B - A_{j_0} y_{j_0} = A_B (A_B^{-1} A_{j_0}) - A_{j_0} = (A_B A_B^{-1}) A_{j_0} - A_{j_0} = A_{j_0} - A_{j_0} = 0.$$

Finally, because

$$\mathbf{c}^T \mathbf{y} = \mathbf{c}_B^T \mathbf{y}_B + c_{j_0} y_{j_0} = \mathbf{c}_B^T (-\mathbf{A}_B^{-1} \mathbf{A}_{j_0}) + c_{j_0} = c_{j_0} - \mathbf{c}_B^T \mathbf{A}_B^{-1} \mathbf{A}_{j_0} = \overline{c}_{j_0}$$

we have that  $\mathbf{c}^T \mathbf{y} = \overline{c}_{j_0} < 0$  which finishes the proof.